

SE3910 – REAL TIME SYSTEMS

Scheduling Theory and Networking

ROADMAP

- Friday
 - Communication between processes
- Monday
 - Gstreamer and OpenCV Introduction
 - Image processing / image handling

Graphics

• Wednesday
⇒ Micro
Measuring Execution
Time.

OBJECTIVES

- Explain rate monotonic analysis and RMA Scheduling
- Explain earliest deadline first scheduling.
- Explain the EDFABound scheduling approach
- Explain the concept of a reader, writer, and bounded buffer
- Explain how Time correlated buffering might be used to aid in the design of a system
- Explain the concept of a ring buffer
- Understand how a mailbox works and why a mailbox might be used.

Finish

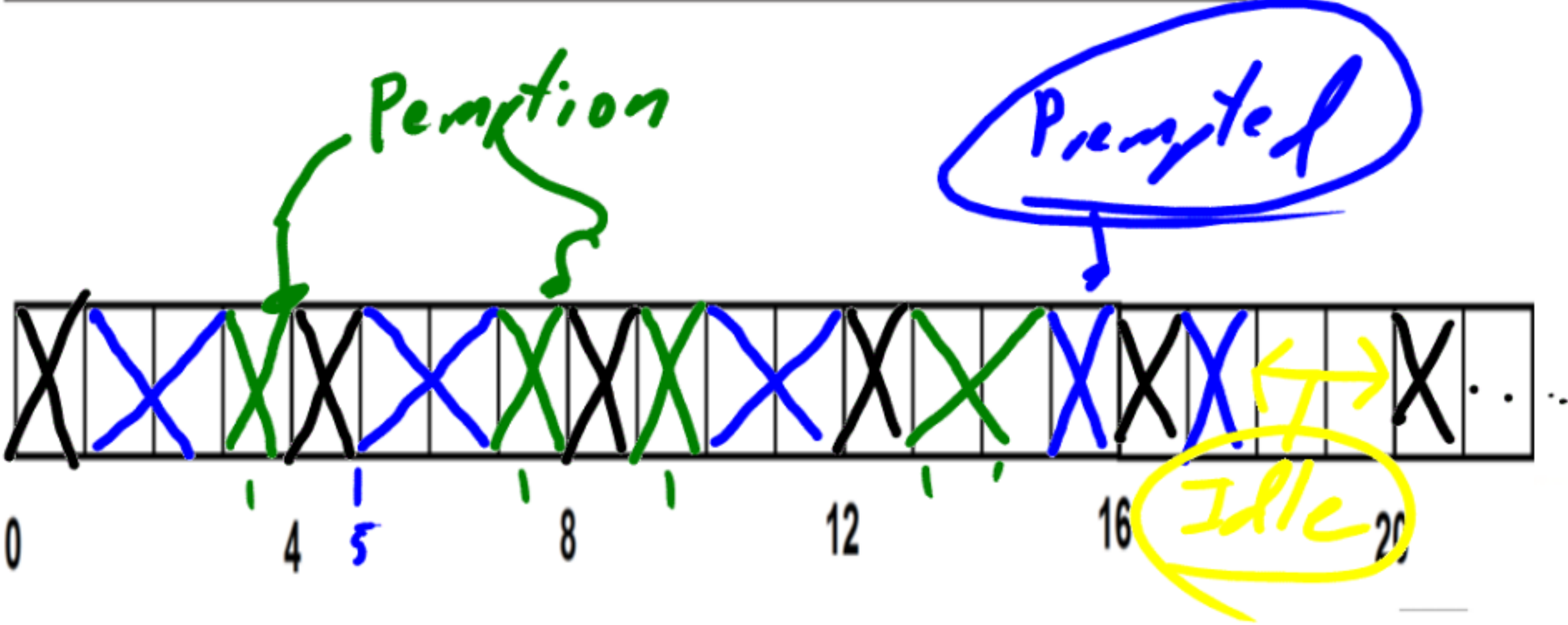
Common

TABLE 3.2. Example Task Set for RM Scheduling

τ_i	p_i	e_i	u_i
τ_1 - Highest	4	1	0.25
τ_2 - Mid	5	2	0.4
τ_3 - Lowest	20	5 ✓	0.25

Processing time

90%



- Max Utilization to be schedulable

$$U \leq n * (2^{\frac{1}{n}} - 1)$$

$$\lim_{n \rightarrow \infty} n * (2^{\frac{1}{n}} - 1) = \ln 2 \approx 0.69$$

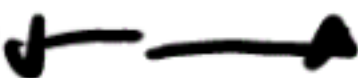
$$3 * 2^{\frac{1}{3}} - 1 = 3 * (1.259 - 1)$$

2.78

RMA BOUND

TABLE 1.3. CPU Utilization (%) Zones

Utilization (%)	Zone Type	Typical Application
<26	Unnecessarily safe	Various
26–50	Very safe	Various
51–68	Safe	Various
69	Theoretical limit	Embedded systems
70–82	Questionable	Embedded systems
83–99	Dangerous	Embedded systems
100	Critical	Marginally stressed systems
>100	Overloaded	Stressed systems



- A dynamic scheduling algorithm
 - Priority of the task changes as other tasks are released and completed.

Theorem: The EDF A Bound

A set of n periodic tasks, each of which relative deadline equals its period, can be feasibly scheduled by the EDF A if and only if $U = \sum_{i=1}^n e_i / p_i \leq 1$.

$$U = \sum_{i=1}^n \frac{e_i}{p_i} \leq 1$$

AN EXAMPLE

TABLE 3.4. Task Pair for the Example of EDF Scheduling

τ_i	p_i	e_i	u_i
τ_1	5	2	0.4 40%
τ_2	7	4	0.57 57%
			<u>97%</u>

- Questions:
 - What is total CPU utilization? **97%**
 - Is this within the RMA bound?

$\Rightarrow 2 \times (2^{\frac{1}{2}} - 1)$

$2 \times (1.414 - 1) = .414 \times 2$

828

TABLE 3.4. Task Pair for the Example of EDF Scheduling

τ_i	p_i	e_i	u_i
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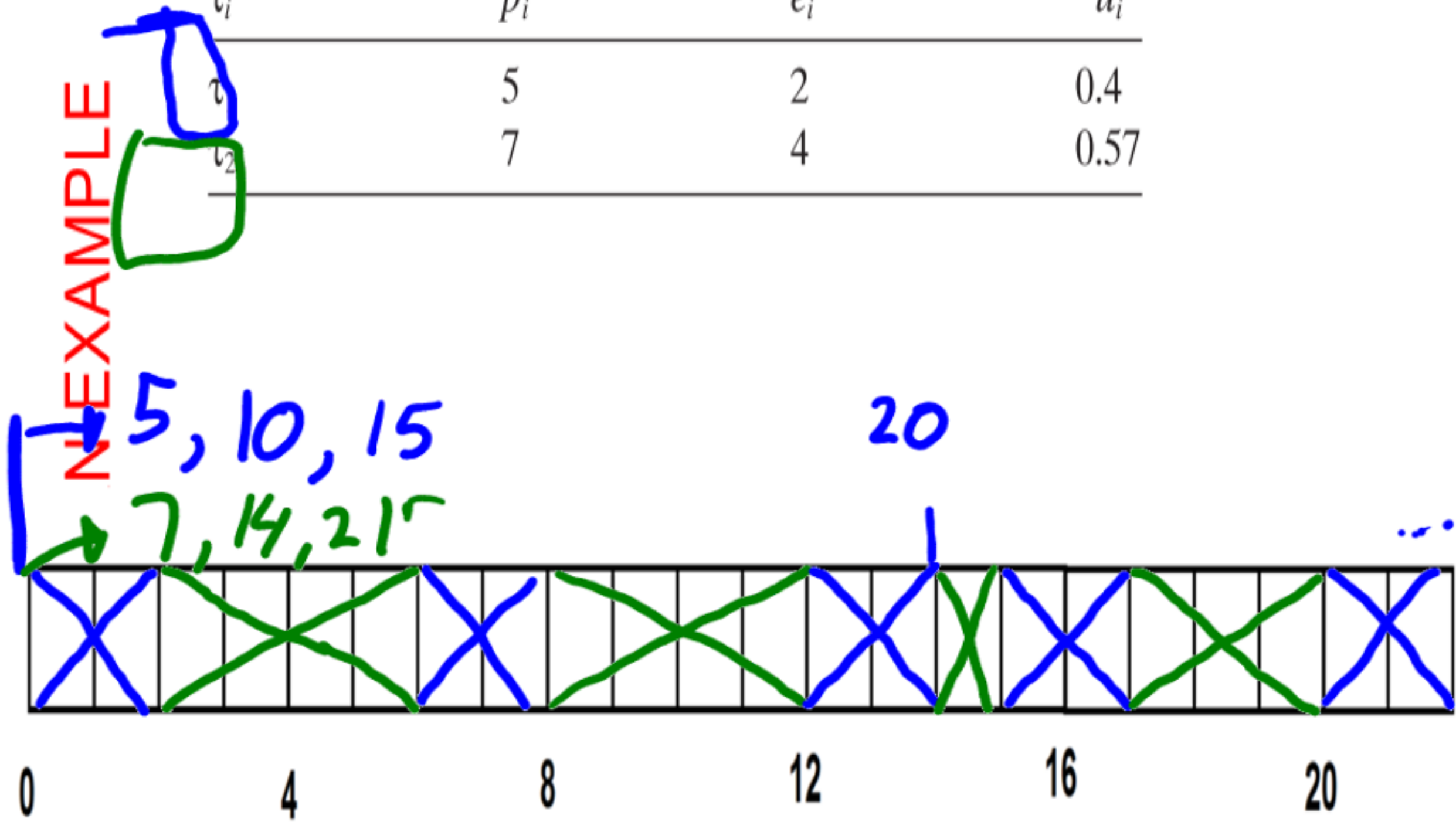


TABLE 3.4. Task Pair for the Example of EDF Scheduling

τ_i	p_i	e_i	u_i
τ_1	5	2	0.4
τ_2	7	4	0.57

N EXAMPLE



0

4

8

12

16

20

ADVANTAGES / DISADVANTAGES

- Dynamic Priority
 - More flexible and achieve higher utilization
- Temporal Behavior \Rightarrow Time
 - More constant with a fixed priority algorithm
 - RMA stable in the presence of missed deadlines
 - Same low priority tasks always miss their deadlines
 - EDF can not predict which tasks will miss their deadlines
- RM tends to use more preemption

COMMUNICATING BETWEEN PROCESSES

- Processes must communicate
 - How can we do that?

*Sending data /
info from
one process
to another.*



COMMUNICATING BETWEEN PROCESSES

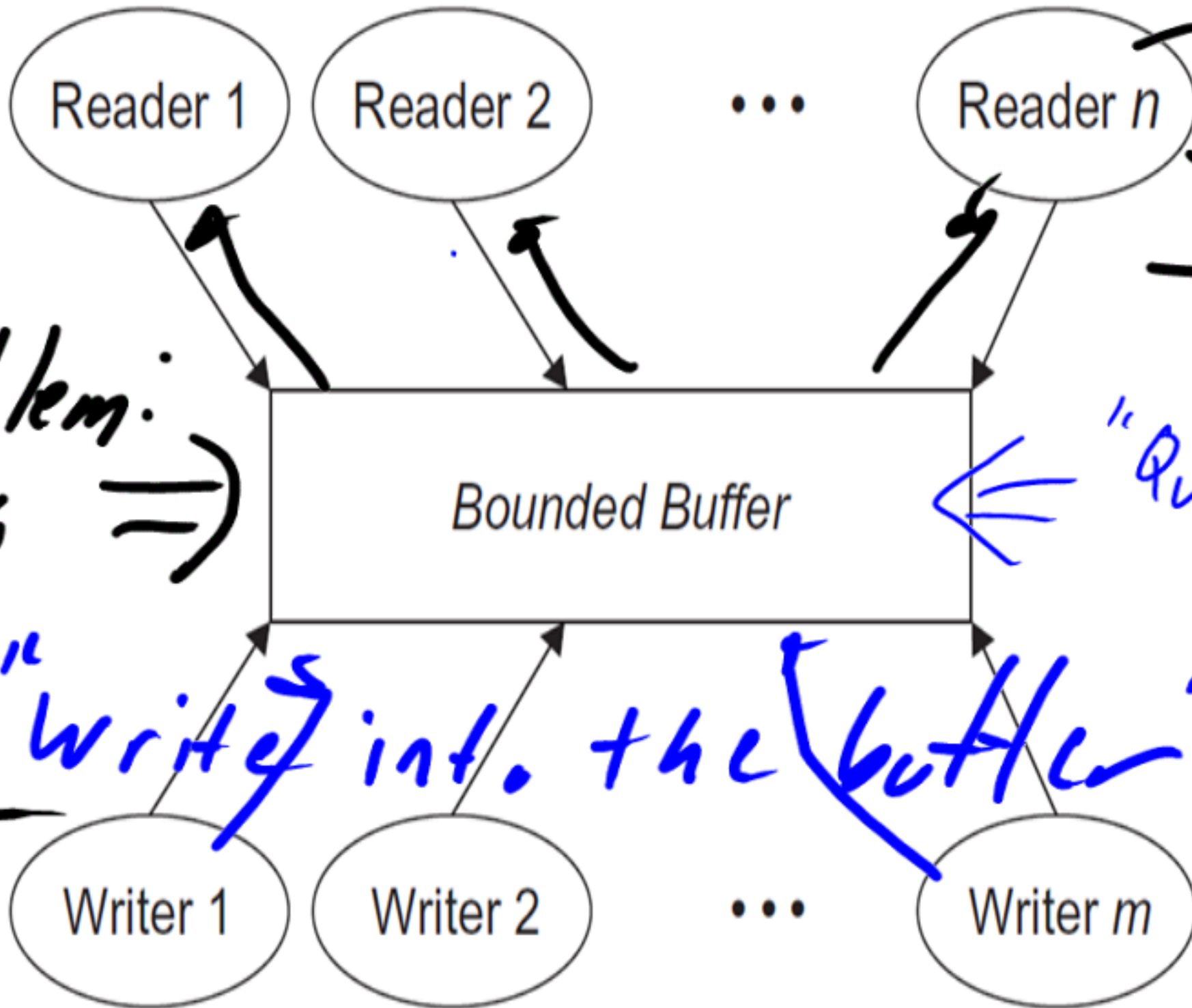
- Processes must communicate
 - How can we do that?

Messages from one process to another.



process

READERS AND WRITERS



Process them

Problem: Lots of data.

"Queue"

"writing into the buffer"

DOUBLE BUFFERING



⇒ "Producer process"



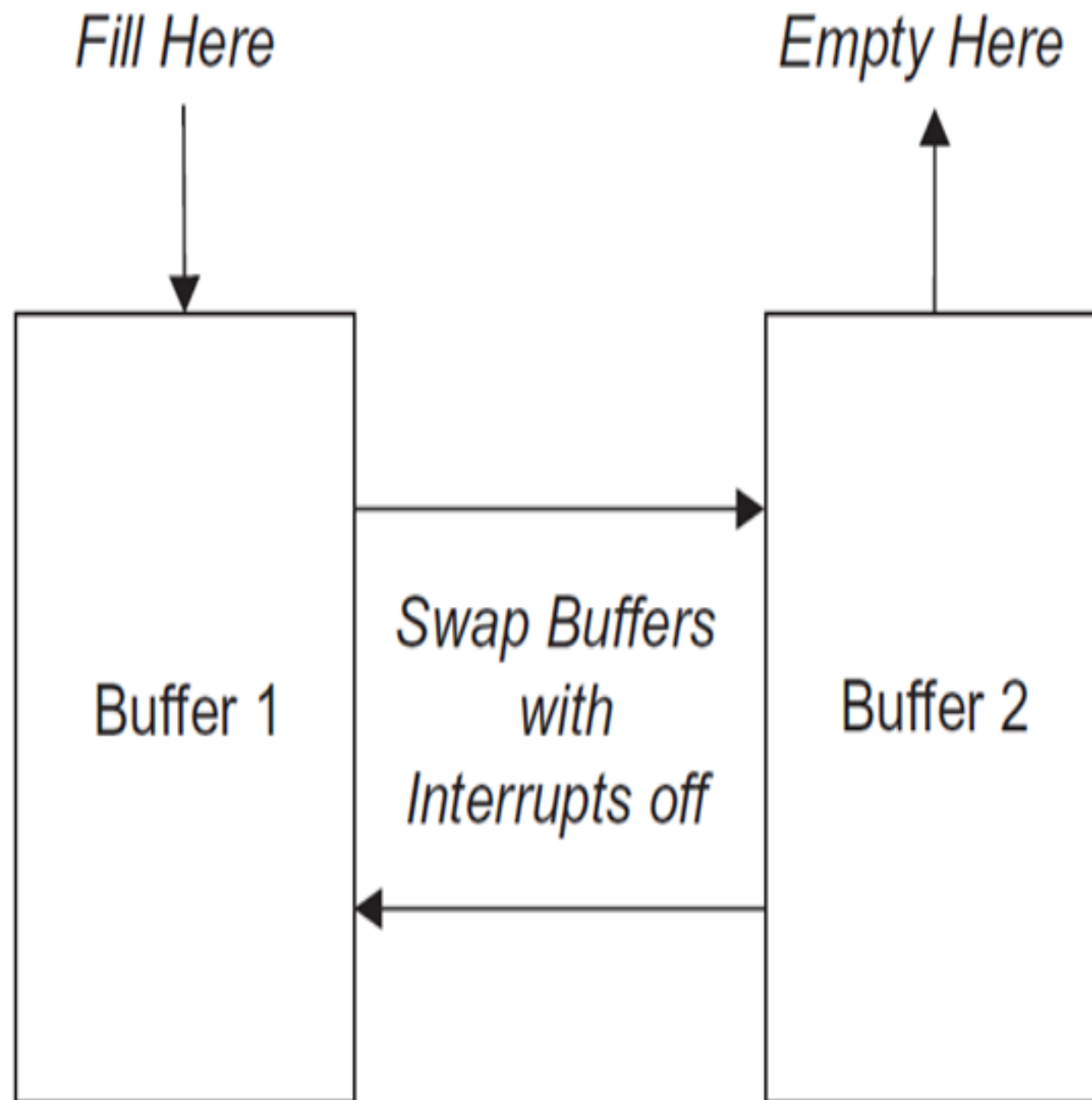
Dump it

Dump



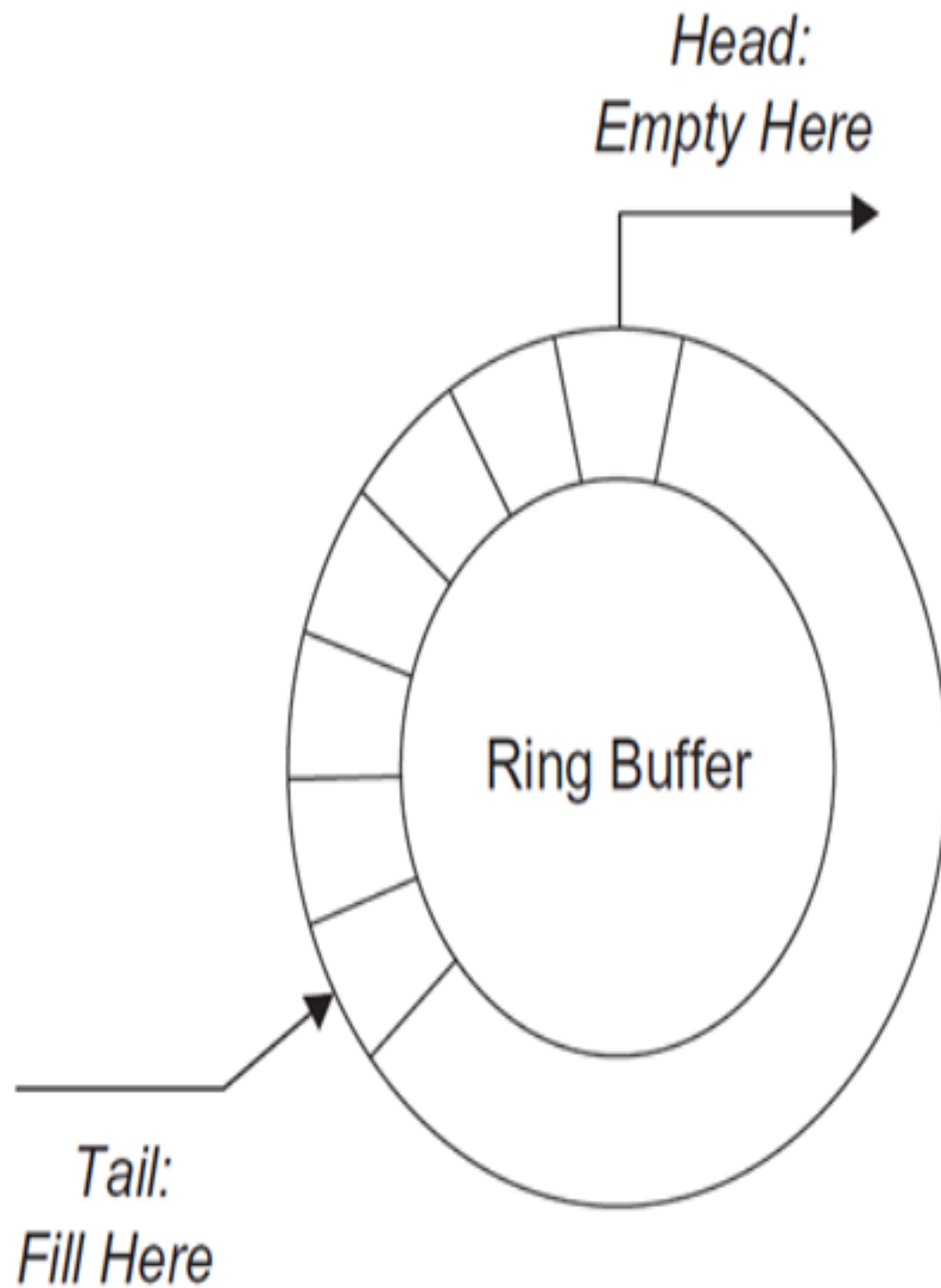
"Consumer process"

DOUBLE BUFFERING DATA STRUCTURE



- We have a GPS System
 - Every 50 ms, we receive a new GPS coordinate
 - 10 ms processing time per coordinate
 - We need to filter the data to remove erroneous readings
 - Filtering requires 5 readings
 - Filtering takes 10 ms
 - After filtering is done, we can refresh the display
 - This takes 100 ms to do.
 - How might we do this?

RING BUFFER



MAILBOXES

- Pend(d, &s)
- Post(d &s)
 - D data that is to be sent / received
 - S The mailbox that is to be referenced.



- Pend
 - Allows multiple processes to use the same mailbox, but only one will receive the material
- May be combined with other data structures to be efficient.