



Instruction Set-Part 1

Lecture Objectives:

- 1) Explain the difference between Harvard and Von Neumann architectures in a computer.
- 2) Define instruction set
- 3) Explain the concept of the source and destination registers for the MIPS instruction set.
- 4) Using the MIPS instruction set, explain how to add a set of variables.
- 5) Define the term computer register
- 6) Define the term data transfer instruction
- 7) Using the MIPS instruction set, initialize a register to a fixed value.

No class Friday

Morning. 

Will be a short video on
website to watch.

Will be a short quiz on

Monday ...

⇒ Over chapter 1 material

⇒ Homework will be e-mailed
out.

- Calculate the geometric mean of the following numbers

$-1, 2, 4$

$n^{\text{th}} \sqrt{\text{Product of } \#}$

$$\sqrt[3]{1 \times 2 \times 4} \Rightarrow \sqrt[3]{8} \Rightarrow 2$$

$-2, 4, 8, 16, 32$

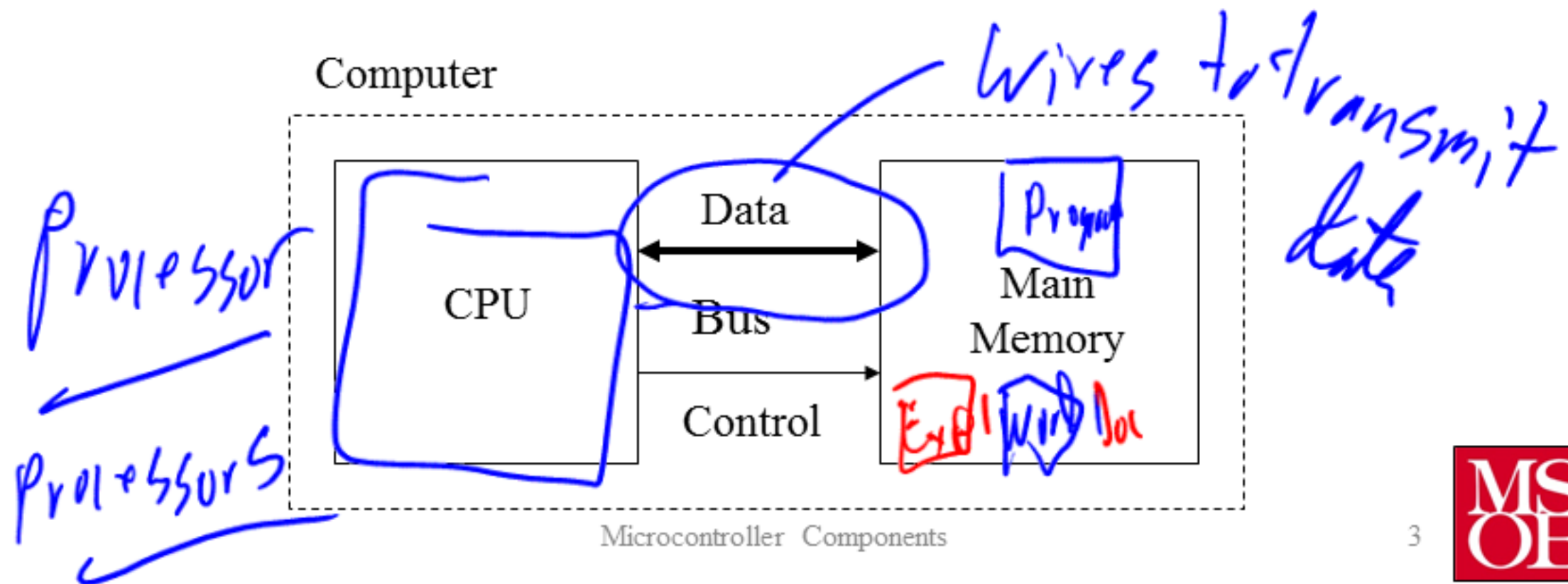
Lead In

$$\sqrt[5]{2 \times 4 \times 8 \times 16 \times 32} \Rightarrow \text{Big } \#$$

$$\sqrt[5]{32768} \Rightarrow 8$$

von Neumann Architecture

- Principles → *More common Architecture*
 - Data and instructions are both stored in the main memory (stored program concept)
 - The content of the memory is addressable by location (without regard to what is stored in that location)
 - Instructions are executed sequentially unless the order is explicitly modified ⇒ *loop, branch, etc.*
 - The basic architecture of the computer consists of:



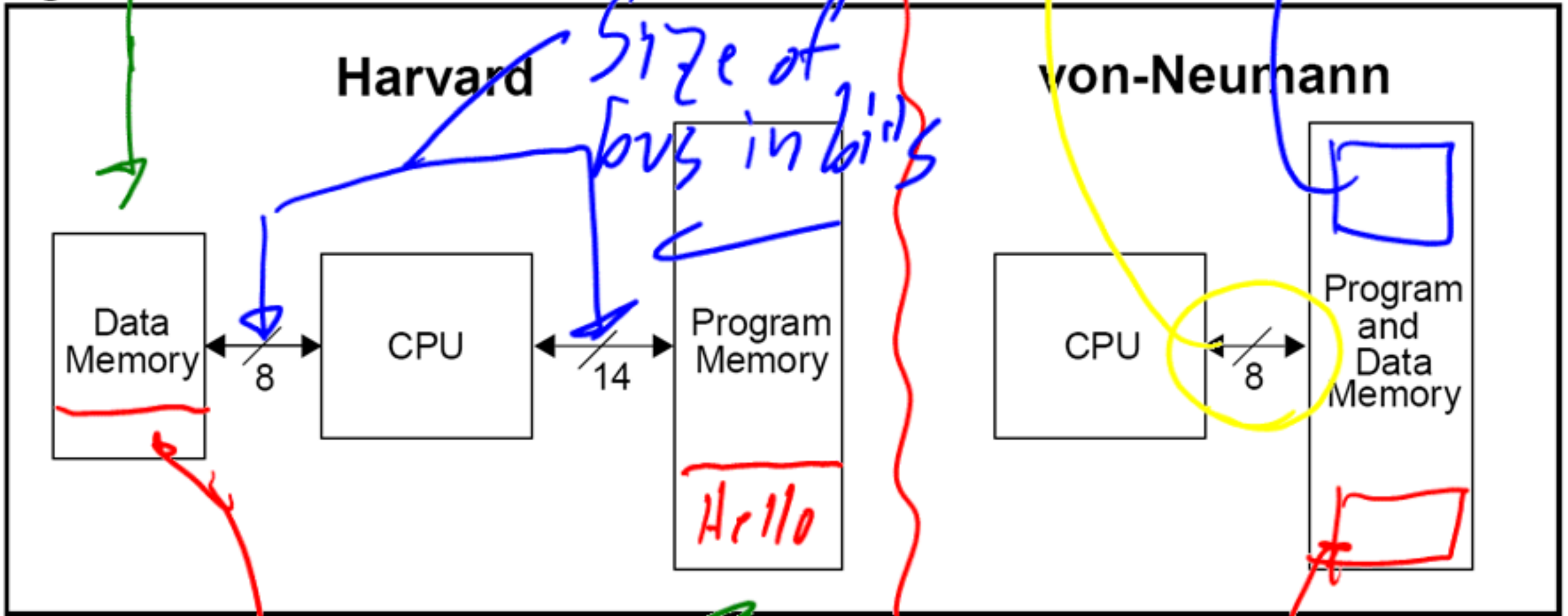
Harvard Architecture

- Assign data and program instructions to different memory spaces. *Programs / Data*
- Each memory space has a separate bus.
- Allows:
 - Different timing, size, and structure for program instructions and data.
 - Concurrent access to data and instructions. *Best for performance*
 - Clear partitioning of data and instructions (=> security)
 - Harder to program

Harvard Versus von Neumann Architecture

Figure

Figure 4-1: Harvard vs. von Neumann Block Architectures



Separate data Memory

64bit's on laptop Hello program

Size of bus in bits

Hello

Hello

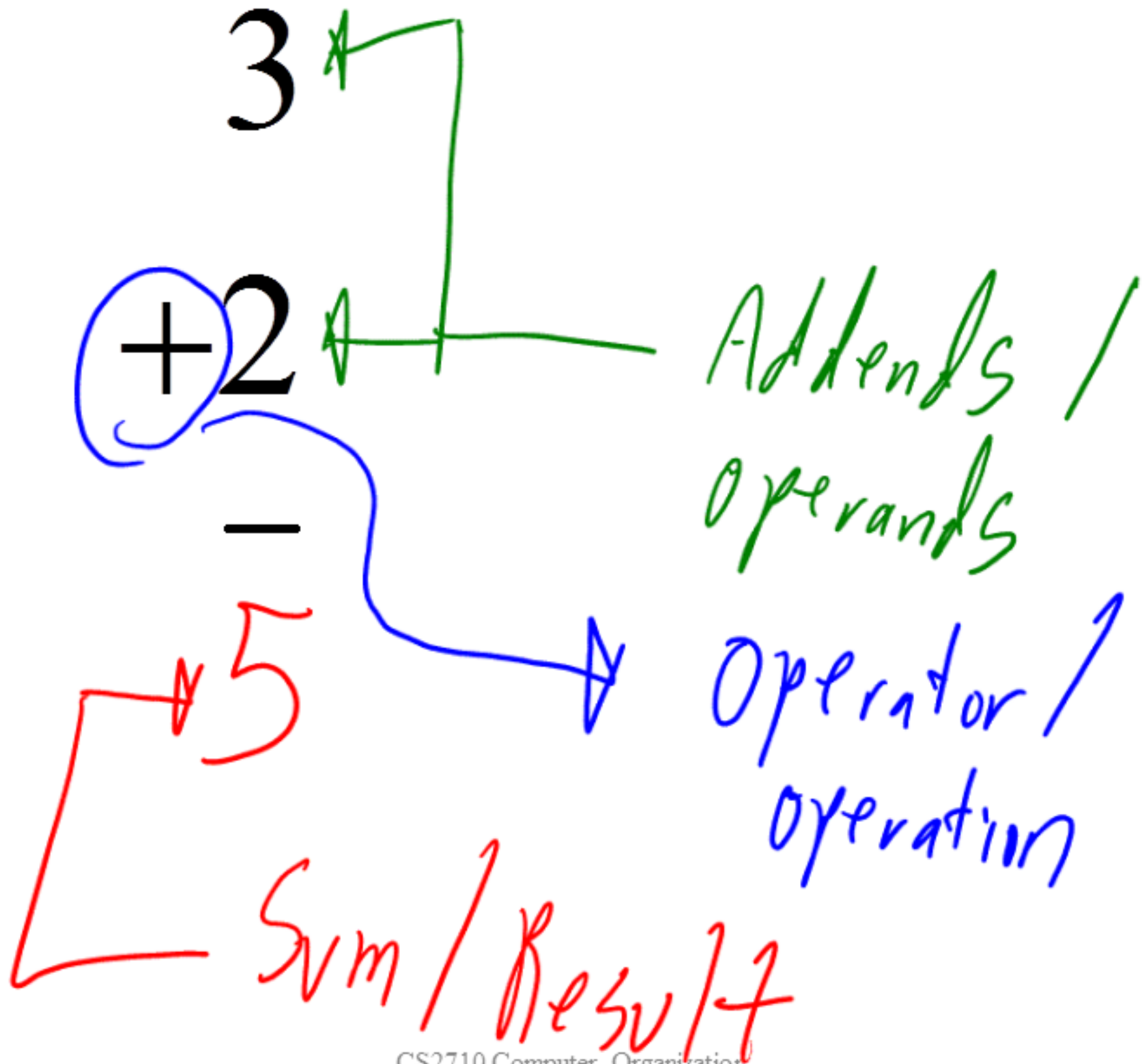
Hi
Hello data Separate Program

Hello data



Operands and Operations in

Math



Mathematics Operations

- Addition -
- Subtraction -
- Multiplication -
- Division -
- Power -
- Square Root -
- Etc. -

Adding

*All those things from
Java.*

Definition of what can
be done on a given computer

Instruction Set

- *The vocabulary of commands understood by a given computer architecture.*
 - The repertoire of instructions of a computer
- Different computers have different instruction sets
 - But with many aspects in common
- Early computers had very simple instruction sets
 - Simplified implementation
- Many modern computers also have simple instruction sets

Stored Program Concept

- The idea that instructions and data of many types can be stored in memory as numbers.

⇒ Hex values



Code ⇒ "Assembly" ⇒ Hex
"code"
Binary code

The MIPS Instruction Set

- Used as the example in textbook *Teaching Language*
- Will be used as example language for the course
- Stanford MIPS commercialized by MIPS Technologies (www.mips.com)
- Large share of embedded core market
 - Applications in consumer electronics, network/storage equipment, cameras, printers, ...
- Typical of many modern ISAs *Instruction Set Architecture*
 - See MIPS Reference Data tear-out card, and Appendixes B and E

Mips is similar to ARM

MIPS Assembly Language

Statement

```
add a, b, c # Adds the sum of b and c and places it in a.
```



Comment

Destination register on
The Mips
Operands
operand

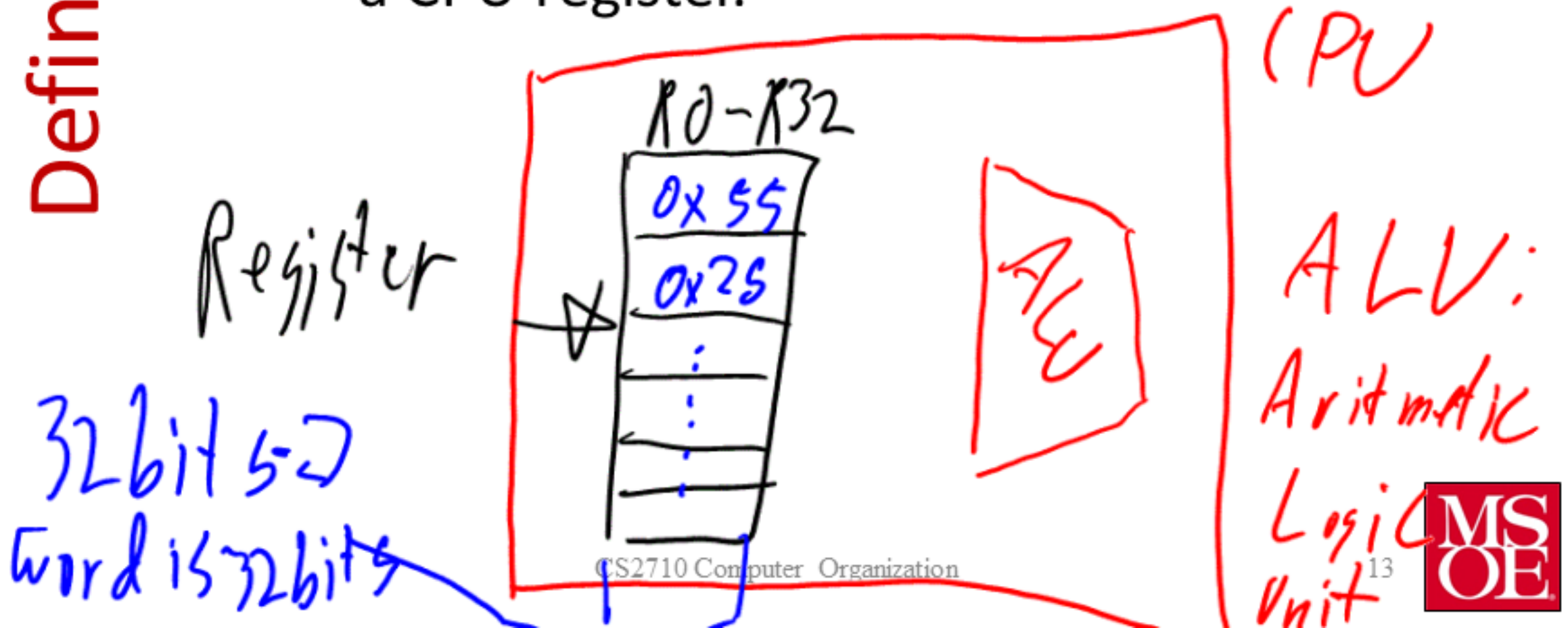
All arithmetic operations have this form.

Design principle 1: Simplicity
- a vor regularity



Definitions

- Register
 - A part of the central processing unit used as a storage location.
- Word *32 bit int in Java*
 - The natural unit of access in a computer.
 - A word normally corresponds to the size of a CPU register.



An Arithmetic Example

- Initial C / Java code:

```
f = (g + h) - (i + j);
```

- Refactored C / Java Code

```
temp0 = g + h;
```

```
temp1 = i + j;
```

```
f = temp0 - temp1;
```

Re-arranged

- MIPS Assembly Code:

```
add t0, g, h # temp t0 = g + h
```

```
add t1, i, j # temp t1 = i + j
```

```
sub f, t0, t1 # f = t0 - t1
```

operands / source registers

t0 / t1 ⇒ Destination Registers

MIPS Architecture

- Arithmetic instructions use register operands
- MIPS has a 32×32 -bit register file
 - Use for frequently accessed data
 - Numbered 0 to 31
 - 32-bit data called a “word”

Fully worked register example

- C code:
 $f = (g + h) - (i + j);$
– f, \dots, j in $\$s0, \dots, \$s4$
- Compiled MIPS code:
add $\$t0, \$s1, \$s2$
add $\$t1, \$s3, \$s4$
sub $\$s0, \$t0, \$t1$

Register Operands

- Assembler names
 - \$t0, \$t1, ..., \$t9 for temporary values
 - \$s0, \$s1, ..., \$s7 for saved variables

Problem

- $r = 6 + 3;$
- How do we do it?